

About New Internal Symmetries of the Integer Lattice (Z^N)

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Abstract

The numerous internal symmetries are found in N-dimensional integer lattices (Z^N). The relation of these symmetries with the new mathematical category named the Masks (or Neighborhoods) is shown. A set of definitions for the Correct Masks and Perfect Masks is presented; the identity between the Correct and Perfect Masks is hypothesized. The relationship between the Perfection of the Mask and the new category named "Mathematical String" is shown. The Correctness of the several Masks in Z^N ($N=1,2$) is proven and a simple method to find the Correctness for all other N is outlined. The hypothesis of high population density of Perfect Masks in integer lattices Z^N with large N is stated.

Keywords: Symmetries; Integer lattice

Introduction

A new set of internal symmetries for N-dimensional integer lattices Z^N are inseparably linked to a new mathematical category named the Mask (Neighborhood). This category represents a successful attempt to generalize the notion of integer number (Z). There exists a generally accepted concept of hierarchy of numbers: Boolean number \Rightarrow natural number \Rightarrow integer number \Rightarrow real number \Rightarrow complex number. The author suggests a slightly different classification: the "tree" forks starting with the integer numbers. After the integer numbers, the standard limb (where all "continuous" numbers belong) goes to the right, and a discrete-space generalization of the "integer number", which is referred in this paper as the Neighborhood (Mask), goes to the left. As the natural numbers can be prime or composite, so the Neighborhoods can be Correct (or Perfect) and... "not quite Correct", i.e., Incorrect. For example: Neighborhood $\{-5;-2;-1;1;2;4\}$ in one dimension is Correct and (apparently) Perfect, yet nothing can be said, for instance, about Neighborhood $\{-5;-3;-1;1;2;4\}$. The notion of Correctness, as we.

Definition of Neighborhood (Mask)

In the beginning, consider replacing the term Neighborhood with its synonym, the Mask. The reasons of doing this will be discussed below.

Let us define Mask M in N-dimensional space as the set of n integer, nonzero, and the different vectors of dimension N:

$$M = \{\overline{a_1}, \overline{a_2}, \dots, \overline{a_n}\}; \overline{a_j} = \{i_1, i_2, \dots, i_N\}; j = 1, \dots, n; i_k \in Z$$

The Mask is normal if among other vectors it comprises all the vectors (we can call it unit vectors) wherein one element i_k ($k=1, \dots, N$) is +1 or -1 while all others elements are 0. The Mask is called primitive if it contains ONLY unit vectors. Thus, primitive Masks are a subclass of normal Masks. The Mask that for any vector $\{i_1, i_2, \dots, i_N\}$ contains vector $\{-i_1, -i_2, \dots, -i_N\}$ is called central-symmetric. In one-dimensional integer space, any Mask is easy to describe with two numbers (m, k) where the binary representation of m gives negative i, and the binary representation of k (after mirroring) gives positive i. For instance, (19,11) stands for Mask $\{-5;-2;-1;1;2;4\}$. This is due to the fact that $(19,11) = (2^4+2^1+2^0, 2^0+2^1+2^3) = (10011; 1101)$; after inverting ("mirroring") the sequence of digits in the second component $(1101)^{Mir} \Rightarrow 1011$, we obtain $\{-5;-2;-1;1;3;4\}$. Similarly, expression $(21,11) = (2^4+2^2+2^0, 2^0+2^1+2^3) = (10101; 1101)^{Mir} = (10101; 1011)$ denotes Mask $\{-5;-3;-1;1;2;4\}$. In the future, we will consider only normal Masks. From this point, any Mask=normal explained in Figure 1.

Definition of the Correct Mask. Theorem 1

Consider lattice Z^N . Some or none of lattice dimensions can be closed (like in a torus) whereas all the other stay infinite, this property has no impact on our future analysis. Take any Mask in N-dimensional integer lattice and add Zero vector to it (all the elements of Zero vector equal 0). Using this Mask with Zero vectors we define a Cellular Automaton (CA; further we refer to this Automaton as Table CA) with the following properties:

1). CA cells can be in one of 6 states. At every moment of time, the state is one of the letters from following set

$$\{-x; -y; -z; +x; +y; +z\}$$

2). The state of each cell in the next moment of time is defined by Transition Tables $R_{-x}, R_{-y}, R_{-z}, R_{+x}, R_{+y}, R_{+z}$:

$$\begin{bmatrix} a_{1,0} & a_{1,0} & \dots & a_{1,0} \\ a_{1,0} & a_{1,0} & \dots & a_{1,0} \\ \dots & \dots & \dots & \dots \\ a_{1,0} & a_{1,0} & \dots & a_{1,0} \end{bmatrix}; a_{i,j} \in \{-x; -y; -z; +x; +y; +z\}$$

The number of columns in the Tables is n (the number points in the Mask) plus 1. Then we take C as the number of rows in the Tables. Both n and C will be constantly used in the text below. Example of Transitional Table for a Mask (1,1); n=2, C=9 (the rule how to find C for a particular Mask will be explained further on):

$$\begin{bmatrix} -y & +y & -y \\ -y & +y & -z \\ -z & +y & -y \\ -y & +y & +y \\ -z & -z & -z \\ -z & -z & +y \\ +y & -z & -z \\ +y & -z & +y \\ \dots & \dots & \dots \end{bmatrix}$$

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Examples of different Masks

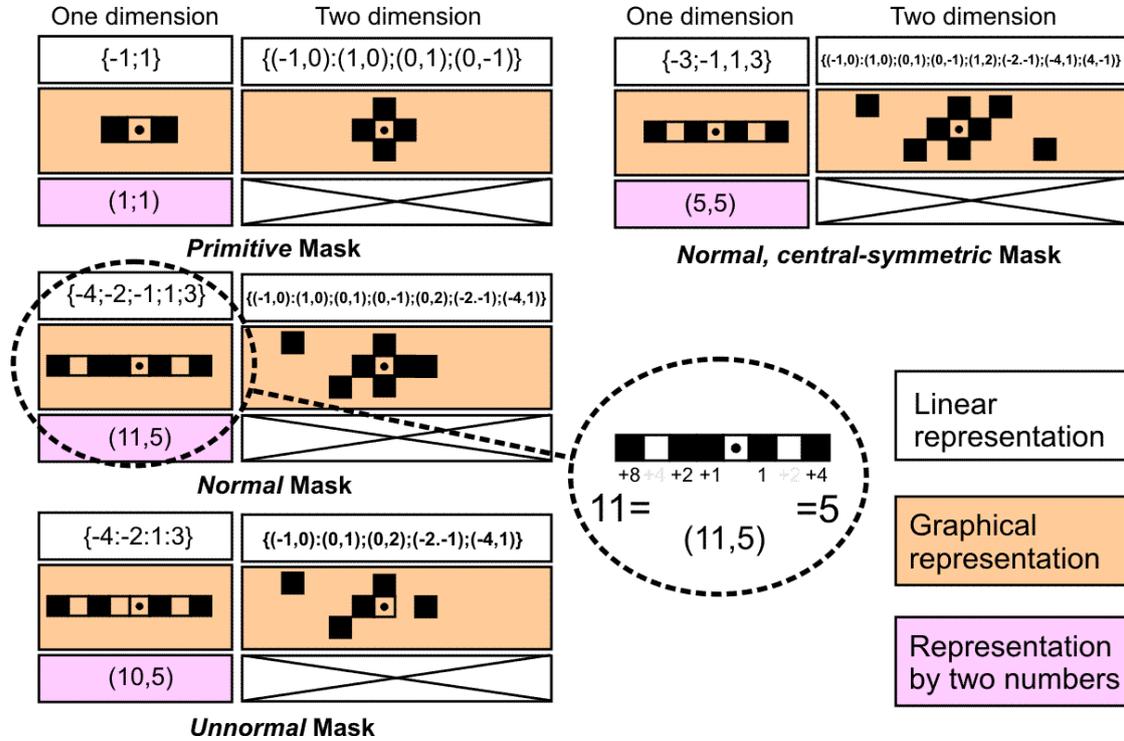


Figure 1: Primitive, normal, and central-symmetric Masks.

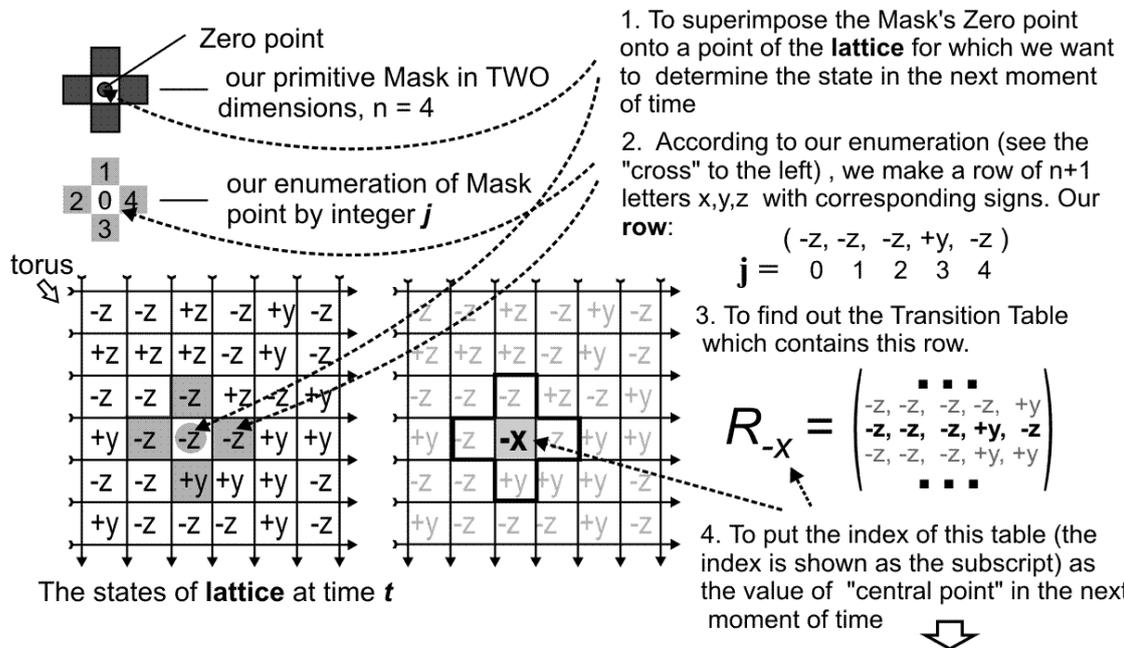


Figure 2: The illustration of one step of the Table CA "run" for a primitive Mask in two dimensions. For clarity, the Automaton run is shown for a single cell.

And similarly for the other letters. We remind here that there are $+1$ columns in Transition Tables R due to inclusion of Zero vector. No two rows of the Transition Tables can be the same. Then a set of standard routines is executed. First, we arbitrarily enumerate our Mask points

plus Zero vector (the center) with numbers j . Figure 2 shows Zero "point" with a dot. Secondly, we superimpose (overlay) each cell in Z^N with zero point of our Mask. Note: it is due to that "overlay" operation, the name "Mask" seems more appropriate than the "Neighborhood".

Thirdly, we write out row $\{a_0, \dots, a_j, \dots, a_n\}$ containing the sequence of states that are taken from the cells on the lattice Z^N where the enumerated vectors comprising our Mask end up illustrated in Figure 2. Fourthly, we check *Transition Tables* $R_x, R_y, R_z, R_{xz}, R_{xy}, R_{yz}$ on a subject of whether or not such a row exists in any of the *Transition Tables*. If we find such a Table, we put its index as the state of CA in the next moment of time ($t+1$). For example, if row $\{a_0, \dots, a_j, \dots, a_n\}$ appears in table R_x , then the next state is $-x$; if some other row appears in table R_y , then the next state is $-y$, etc. Because no rows of the *Transition Tables* can be the same, there is only one state of lattice Z^N in each time point. Figure 2 shows the transformation of arbitrary selected element from integer lattice Z^2 . It illustrates a single step of operation of a *Table CA* for a *primitive Mask* in two dimensions on a closed lattice (i.e., the lattice without "boundary") in this particular case, a surface of 6×6 torus. In accordance to CA definition, the same operation should be performed at once for all cells of CA Figure 3. Let us compare, for example, the run of our Cellular Automaton with that of well-known CA named *Rule 30* by Stephen Wolfram [1] (Figure 4).

Rule 30 has 2 states of the cells while our CA has 6. *Rule 30* works or runs (if we consider the "works" through our terms) for a *primitive Mask* (1,1) in one dimension. The binary decomposition of 30 is abbreviation for the transition rules [1] and Table 1. From those rules, we can determine the *Transition Tables*:

Current Pattern	111	110	101	100	011	010	001	000
New state for center cell	0	0	0	1	1	1	1	0

Table 1: Transition rules for Rule 30 Automaton.

$$R_0 = \begin{bmatrix} 0 & 0 & 0 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}; R_1 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix} \tag{4}$$

We can notice that contrary to our *Automaton*, *CA Rule 30* is not reversible (see explanations below). Any *Transition Table* has a dedicated column (numbered as "h") which is associated with the Zero Vector that we added to the Mask. It is logical to assume $h=0$ for the dimensions greater than one. Yet, in the case of one dimension, it is more convenient to leave the "natural" order (for instance, $h=1$ for *Automaton Rule 30*).

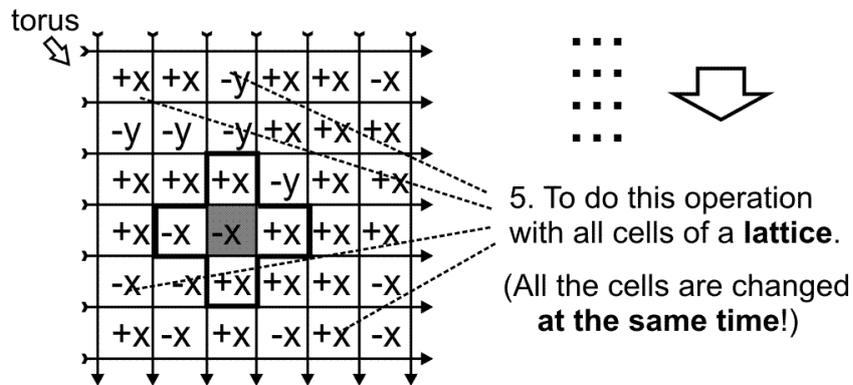
3) The obtained *Transition Tables* are interdependent and can be obtained from table R_x after following substitutions:

$$R_x: e = \{-x, -y, -z, +x, +y, +z\}; R_y: e_1 = \{-z, -x, -y, +z, +x, +y\}; R_z: e_2 = \{-y, -z, -x, +y, +z, +x\}; \tag{5}$$

$$R_{xz}: e = \{+x, +z, +y, -x, -z, -y\}; R_{xy}: e_4 = \{+y, +x, +z, -y, -x, -z\}; R_{yz}: e_5 = \{z, +y, +x, -z, -y, -x\}$$

These substitutions form a group; its multiplication table ($u \times v$) is shown in Figure 5.

4) Our *Automaton* is reversible. The *inverting Transition Tables* are calculated from the same Table.



The states of a lattice at time $t + 1$

Figure 3: The operation shown in Figure 2 are performed for all the cells at once.

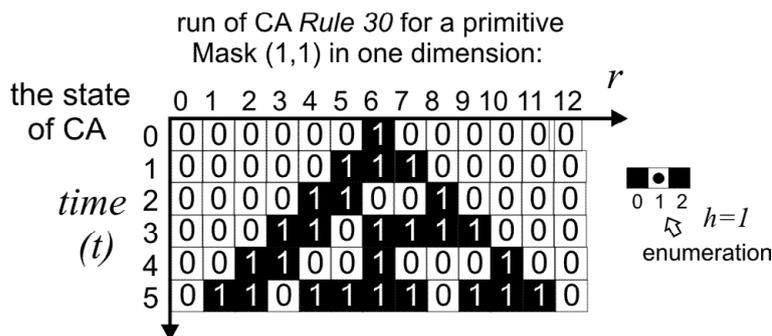


Figure 4: The plot "Automaton state vs. Time" for Wolfram's Automaton "30".

R_x using following substitutions:

$$R_{-x}^{-1} : \{-x, -z, -y, +z, +y, +x\}; R_{-y}^{-1} : \{-y, -x, -z + x, +z, +y\}; R_{-z}^{-1} : \{-z, -y, -x + y, +x, +z\};$$

$$R_{+x}^{-1} : \{+z, +x, +y, -x, -y, -z\}; R_{+y}^{-1} : \{+y, +z, +x, -z, -x, -y\}; R_{+z}^{-1} : \{+x, +y, +z, -y, -z, -x\}.$$

5) When getting started from any initial state of the integer lattice (\mathbb{Z}^N) filled with only $-x$ and $+x$ (further we refer to this the state as the Beginning Point or BP), our CA will last forever. In other words, all the time as CA runs, only those states appear which are present in tables R.

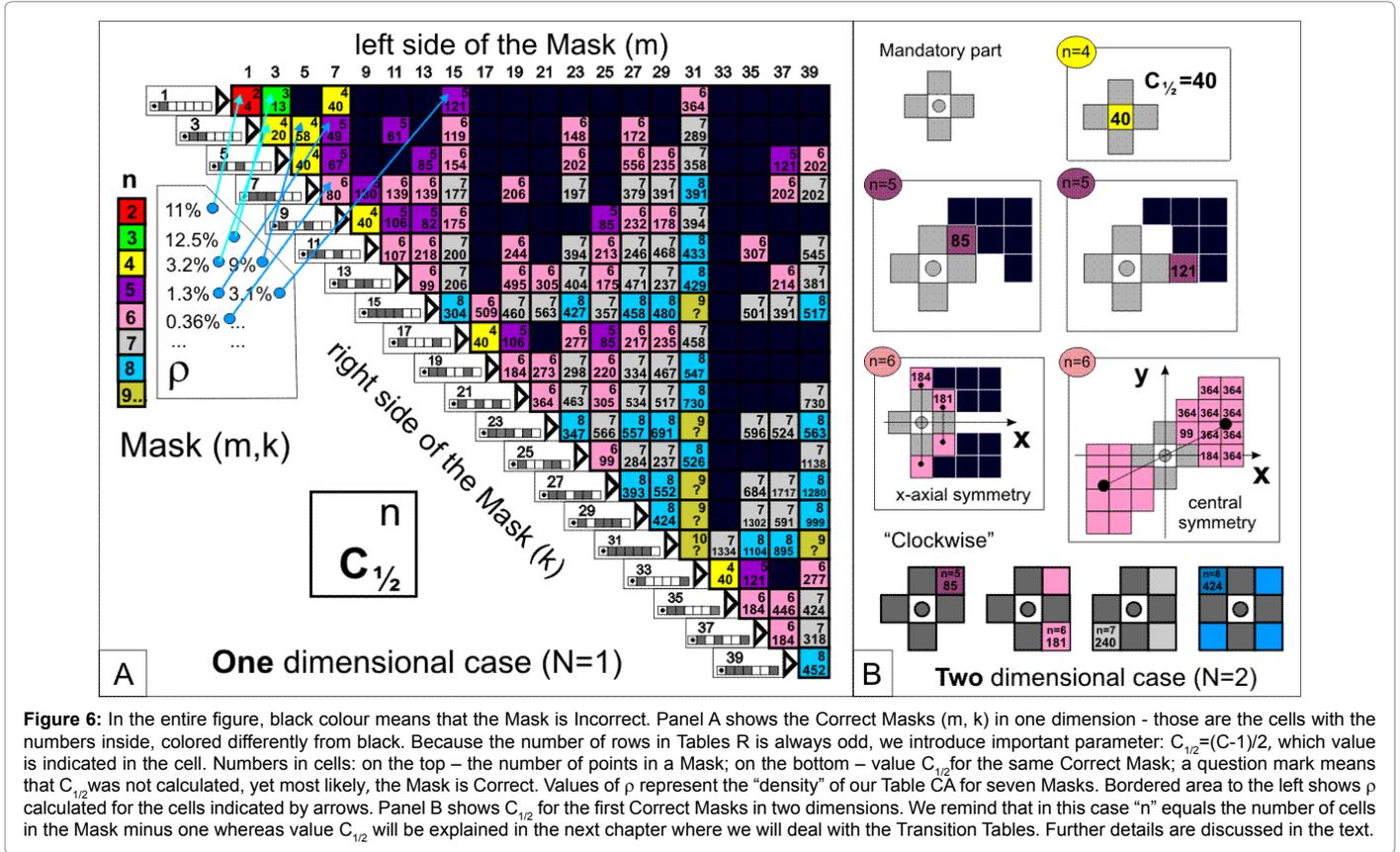
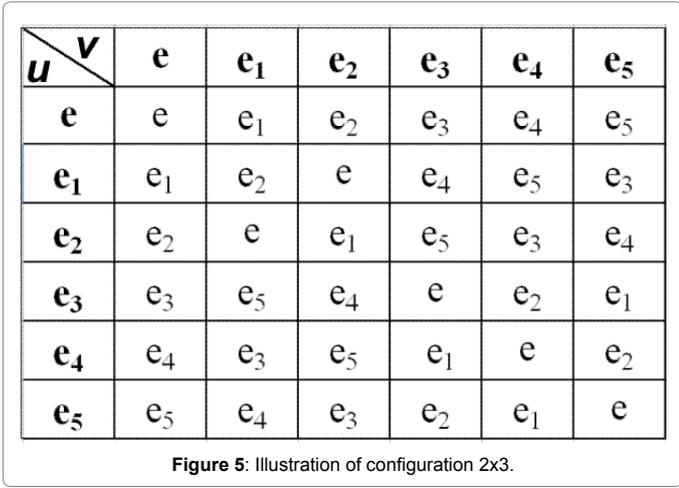
6) At least one row containing only letters $-x$ and $+x$ must be found in tables R_y or R_x or R_z

or R_{+z} . (In other words, the "unionized" table $R_y \cup R_{+y} \cup R_{-z} \cup R_{+z}$ should contain this row).

7) All points (vectors) in the Mask are important. In other words, no point of the Mask can be safely eliminated without changing the Automaton.

If, for a given Mask, it is possible to find CA meeting all seven conditions, then we call this Mask *Correct*. If it is not possible (or until proven otherwise), we call such a Mask *Incorrect*. Let us explain why the very existence of *Correct* Masks is a "miracle" of its kind. It is far from being obvious that at least one Mask is *Correct* in some N-dimensional space. From point 5) it follows that the Automaton must be able to make at least the first step, therefore, the Tables should contain all the 2^{n+1} states formed from the letters of the same type " x ". At some initial conditions after the first step (see point 6) there must appear a new letter (not $+x$ or $-x$) in the states. Now, we are curious about the way the rows containing these letters are transformed. Assuming that we know the answer, it will immediately be transferred into six *Transition Tables* by performing substitutions, see expressions (5). Note that it will also be transferred into six *Inverting Tables* and, in the end, into the same *Transition Tables*. Each step causes new transfers. It may very soon turn out that two identical rows appear in two different tables, and this is prohibited. Yet, a situation is not that bad. As a matter of fact, there is a huge number of the *Correct* Masks and apparently in all dimensions.

Theorem 1. All Masks in one and two dimensions which are not colored by black in Figure 6 (at least with $n < 9$) are *Correct*. Panel B in Figure 6 also shows: the upper row – «a mandatory part» of all *normal* Masks in two dimensions and a *primitive* Mask ("cross" with $n=4, C_{1/2}=40$); the second row – examples of Masks for $n=5$; the third



row – examples of Masks for $n=6$; and the bottom row – the clockwise “attachment” of the cells to the *primitive* Mask that produces the *Correct* Masks with $n=5, 6, 7$, and 8 . It is “obvious” that all *primitive* Masks in all dimensions are *Correct*, and values C for these Masks equal 9^{N-1} for any N , where N is the dimension of the space. The second row (Figure 6B) shows two examples of *Correct* Masks for $n=5$ composed by the addition of points $(x=1, y=1)$ or $(x=2, y=0)$. The third row of panel B in Figure 6 also shows *Correct* Masks for $n=6$: one way to build the *Correct* Mask for $n=6$ is using the x -axis symmetry and another one, using central symmetry ($C_{1/2}$ for $n=6$ are shown in the corresponding cells). In the bottom panel, we add the cells, one by one, to the mandatory part of the two-dimensional Mask starting from the top-right “cell” and then going clockwise. We obtain four *Correct* Masks: $(n=5, C_{1/2}=85)$; $(n=6, C_{1/2}=181)$; $(n=7, C_{1/2}=240)$; and $(n=8, C_{1/2}=424)$. The black cells in Figure 6B indicate some “additions” that yield the *Incorrect* Mask. Parameter p in Figure 6 is the “density” of corresponding CA, e.g., 11% for Mask (1,1); 12.5 % for Mask (3,1) etc. We define here the “density” as the ratio of the number of rows in six matrices R to a number of all possible states $6C/6^{n+1}=C/6^n$ expressed in percent. It is hard to imagine

how accurately our CA must walk! In all states, except the first two, more than 90% of the rows are “forbidden”. Yet, those CA in its time propagation never set a foot on the “forbidden” ground. That is why it can be considered as a “miracle” of its kind.

Proof: Note, the proof for Theorem 1 is obtained using personal computer. The software program with the source code (common to dimensions $N=1$ and $N=2$). In the beginning, we present the *Transition Tables* for *primitive* Masks in one (Figure 7A and 7B) dimensions as well as for Mask (3,1) in one dimension (Figure 7C). Let us explain what is depicted in Figure 7. In each of three panels, we move from the left to the right. Some rows in Figure 7 containing only letter “y” and marked with inclined lines (shading) indicate that these rows are important to perform the “mathematical induction” step in our proof, which is conducted further on in the text. Let us consider letter “y” (note that in case of letter “x” everything is considered similarly). To prove the *Correctness* of our Mask, first, we need to make sure that our table contains all the 2^{n+1} rows consisting of one and the same letter “x”. The R_x table has a row containing only “-z”. Therefore, table R_z

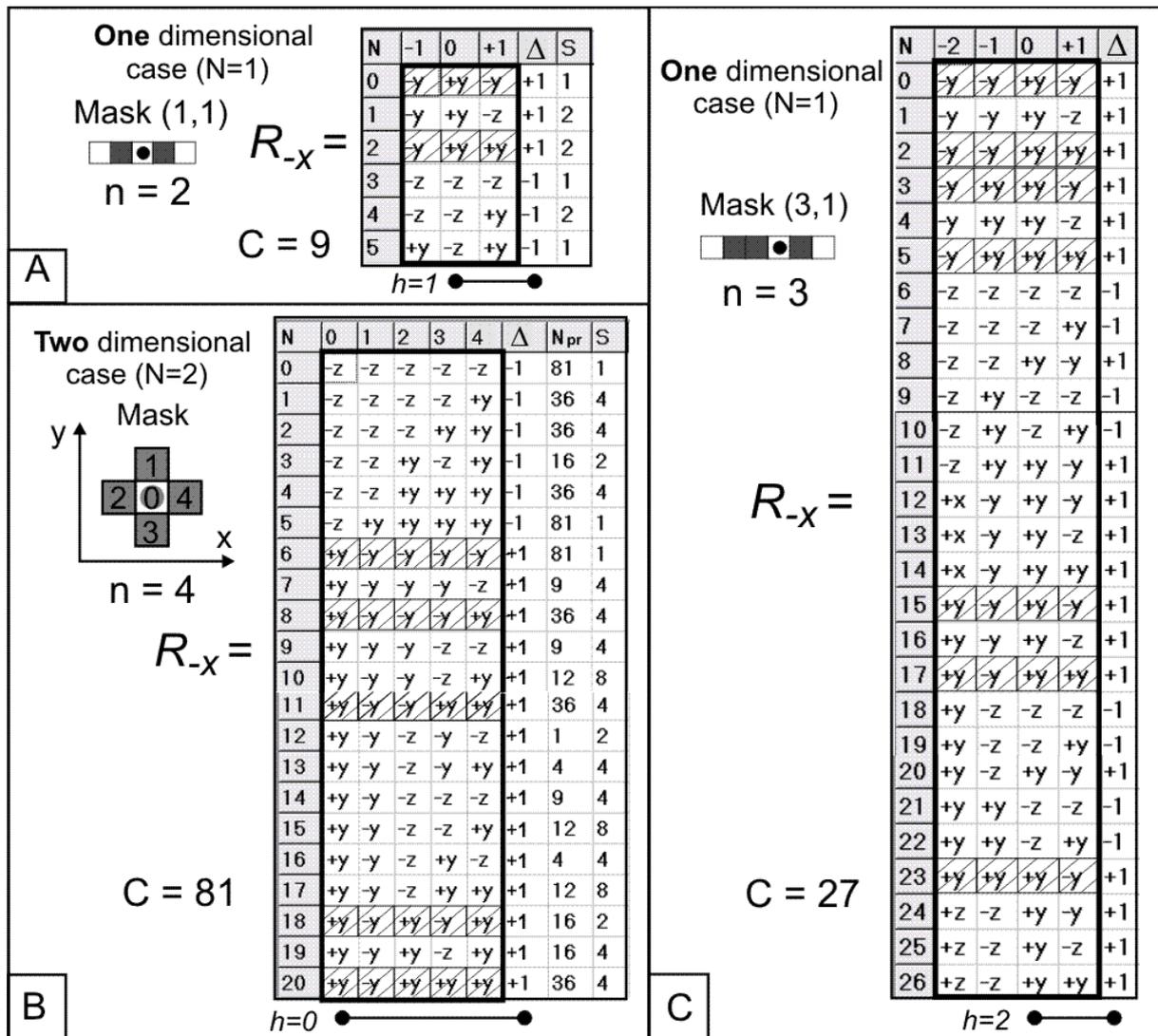


Figure 7: Tables R_x for *Correct* Masks. Panel A for Mask (1,1); panel B for *primitive* Mask in two dimensions; panel C for Mask (3,1) in one dimension. The explanation is in the text.

has a row containing only “-y” (see point 4), and table R_{+x} has a row containing only “+y”. The remaining $(2^{n+1}-2)/2$ variants are distributed over tables R_x and R_{+z} . Moreover, the variants in R_x are complementary to variants R_{+z} (note that all rows from R_x are multiplied by -1 in R_{+z}). Let us check the number of shaded rows. Consider panel A (Figure 7): on one hand, we have $(2^{2+1}-2)/2=3$ rows containing only letter “y”. On the other hand, if in the same panel A, we sum up the numbers in the column S for rows containing only “y” (S stands for symmetry, explained below) then $1+2=3$. Panel B (Figure 7): similarly, using the above expression we can find $(2^{3+1}-2)/2=7$ rows or, on another hand, summing up number in columns S, we get the same number 7 (S=1 everywhere because of no symmetry for the Mask). Finally, panel C: $(2^{4+1}-2)/2=15$ or summing up the rows with account of their symmetry we get $1+4+4+2+4=15$ lines.

Column “Δ” is just a sign of a letter in “zero” point (column) of table R (the relationship is shown by the connected dots; Figure 7). More detailed information about that column will be given in the next chapter. Column “ N_{pr} ” (Figure 7C) shows the number of instances of the given configuration appeared in our proof (see below). The values in this column are given in multiples of C (i.e., this entire column should be multiplied by 81). Then for Figure.7A and 7B it turns out that this entire column is filled only with value 1. We could have hypothesized that all the values in column N_{pr} were multiples of C all the time. But this is not so. This hypothesis turns invalid already for the next Correct Mask (5,3). Column “S” serves to reduce the list of entries. Each row can be found S times in that list, taking into account the axial symmetry (Figure 7A) or square symmetry (Figure 7B).

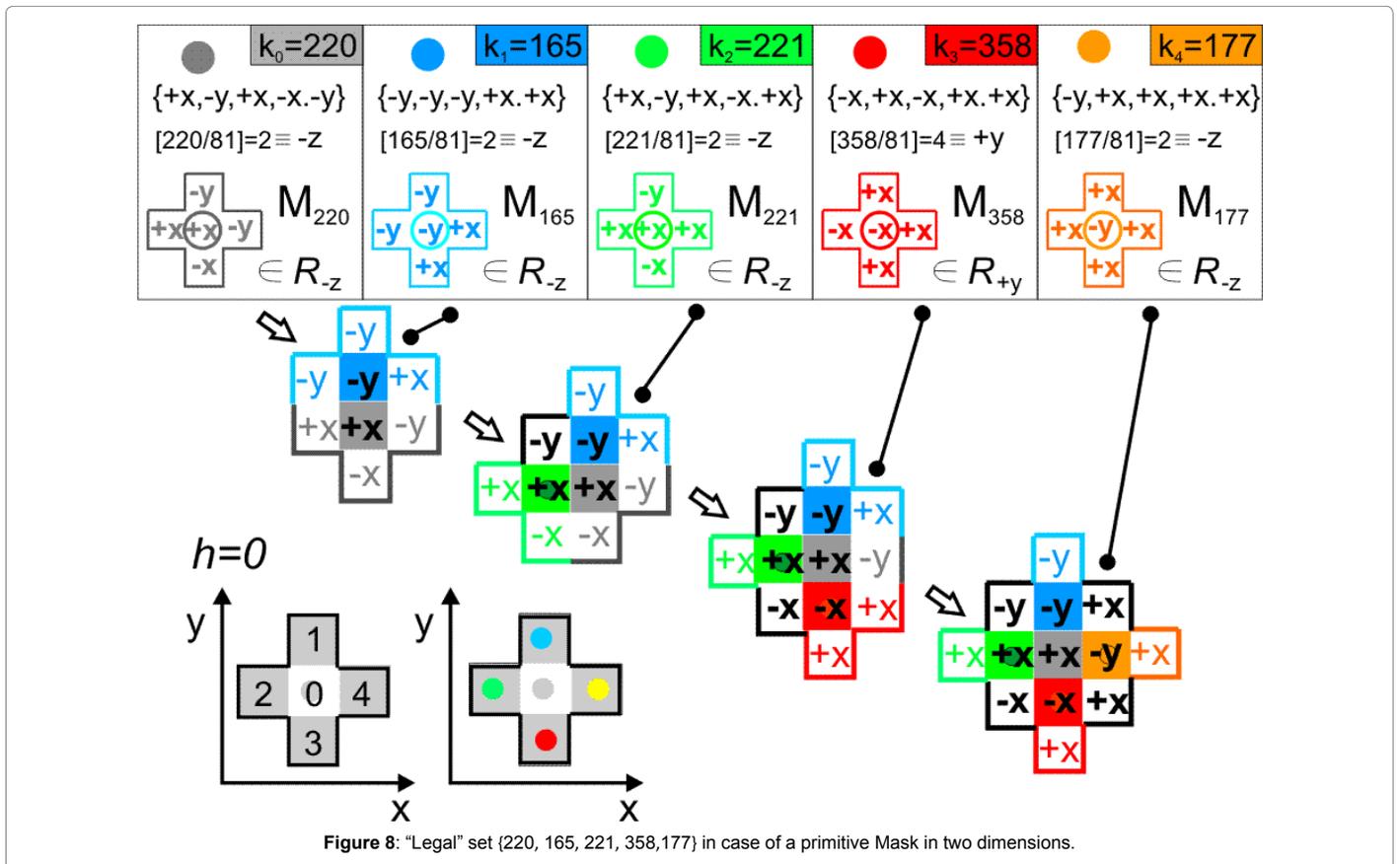
Let us return to the proof. If table R_x is already known for some

Mask, then the following action on the proof is clear. Yet, we are not certain about the outcome of such a process.

We made the first step in the mathematical induction, i.e., we showed that all 2^{n+1} rows containing only letter “x” are present in the Transition Tables. Now, let us move on. Given: assume that at time point t in every point of CA the content of the Mask is permitted, i.e. all the rows in the Masks belong to tables R. It is necessary to prove that at the next time point $(t+1)$ our CA will have the same property. The solution can be found through the proof by exhaustion performed on a personal computer.

Let us unite all Transitions Tables ($R_x, R_y, R_z, R_{+x}, R_{+y}, R_{+z}$) into one Table and named it $R^{(6)}$. It has 6C rows numbered 0, 1, ... k, ... 6C-1 and the numbering of its rows has the following construction. First, the rows of table R_x are placed in a descending order in accordance with the selected seniority - x>-y>-z>+x>+y>+z. Other Transition Tables can be found by using corresponding substitutions, i.e., using the rule (5) from point 3, and numbering in these Tables does not change relatively to that in R_x . After that, the Tables are concatenated (here it means the second one is attached to the bottom of the first one, etc) in similar order: $R \dots, R_y \dots R_z \dots R_{+x} \dots R_{+y} \dots R_{+z}$. To determine the specific Table for the row with number k, one has to determine the floor function $[k/C]$, where C is the number of rows in a single Table of our Mask.

Let us introduce a new term. Let us name M_{kj} as the Mask filled with the row numbered kj from $R^{(6)}$; $k_j=0, \dots, 6C-1$ (see an example of such a filling in Figure 8). We call set $\{k_0, k_1, \dots, k_n\}$ from n+1 (n-number of the points in the Mask) filled Masks “illegal” if after the placement of set M_{k_0} into a point with number 0 of our Masks, M_{k_1} - into



a point with number of 1 our Mask, etc., some point of our lattice will be filled with two different letters, that is the letters not equal to each other. Otherwise we call our set "legal". Now consider, for example, the primitive Mask in two dimensions (see panel B in Figure 7). In Figure 7 (panel B), a natural sequence of row numbering in the considered Transition Table $R^{(6)}$ is truncated (or skipped) in many places due to symmetries (e.g., instead of writing in three rows, one after another, we have placed a single row and indicated the corresponding symmetry). Yet, if we write out all the rows of combined Transition Table $R^{(6)}$ as they are, then it will turn out, that for instance, rows number $k_0=220; k_1=165; k_2=221; k_3=358; k_4=177$ form a "legal" set (Figure 8). It is clear that only a limited number of all sets are "legal". For example, it is easy to see, that if we have any "legal" combination $\{k_0, k_1, \dots, k_n\}$ and $h=0$ (Figure 7 to determine h) then any set $\{m, k_1, \dots, k_n\}$ where m it is not equal k_0 , will be "illegal". Further we find "legal" sets $\{k_0, k_1, \dots, k_n\}$ using the $n+1$ nested loops (each loop considers 6C variants for k_i) thus going through all possible combinations of k_i (Figure 9).

For each "legal" combination, we determine a row in the Mask in consecutive point in time $\{[k_0/C], [k_1/C], \dots, [k_n/C]\}$ and make two checks: i) whether or not the row under consideration can be found in $R^{(6)}$ and ii) whether or not the inverse transformation works. After that, we add "one" to column N_{pr} corresponding to that row and move on to a next step in a cycle... We reach the end, and declare that the Theorem is proved. So, we have got a key observation. If we have the Transition Tables the proof of Correctness is always leads to a success. This fact can be considered as the greatest mystery in mathematics.

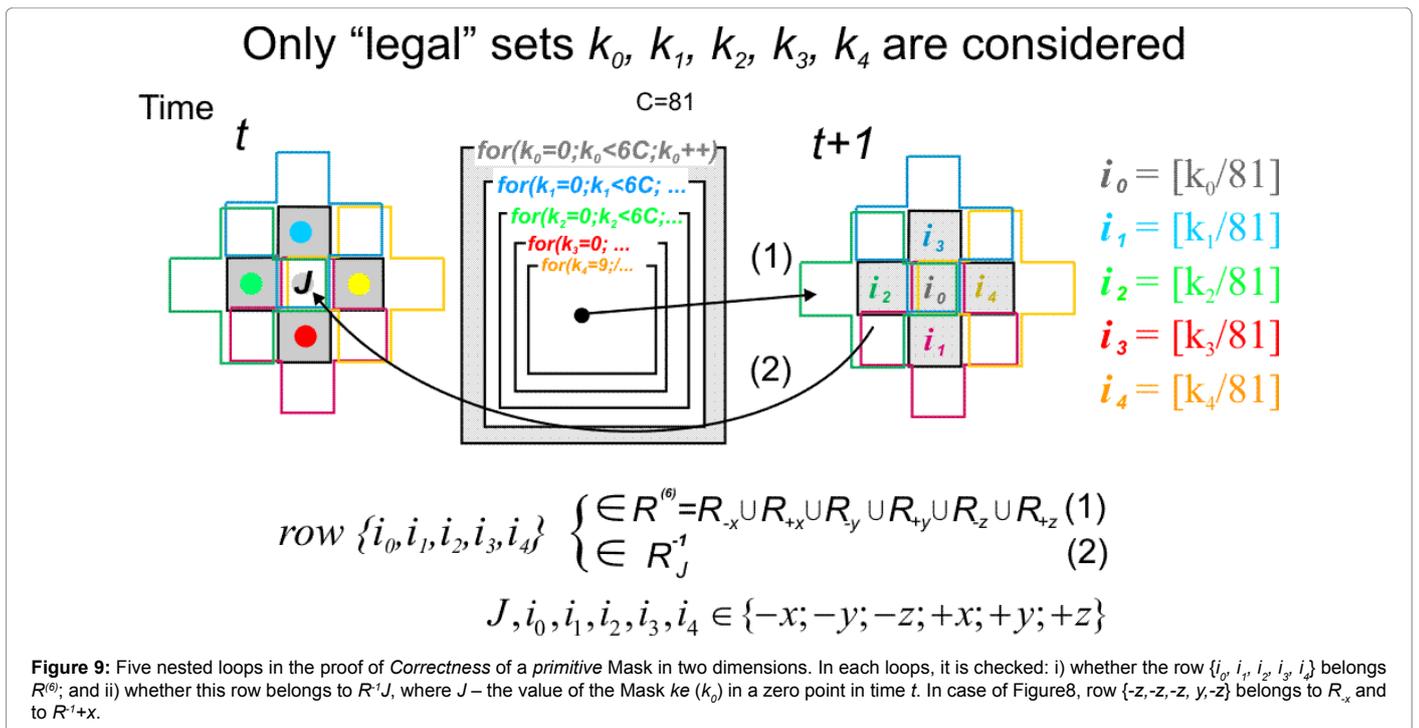
Definition of Simple reversible Automaton for the Mask

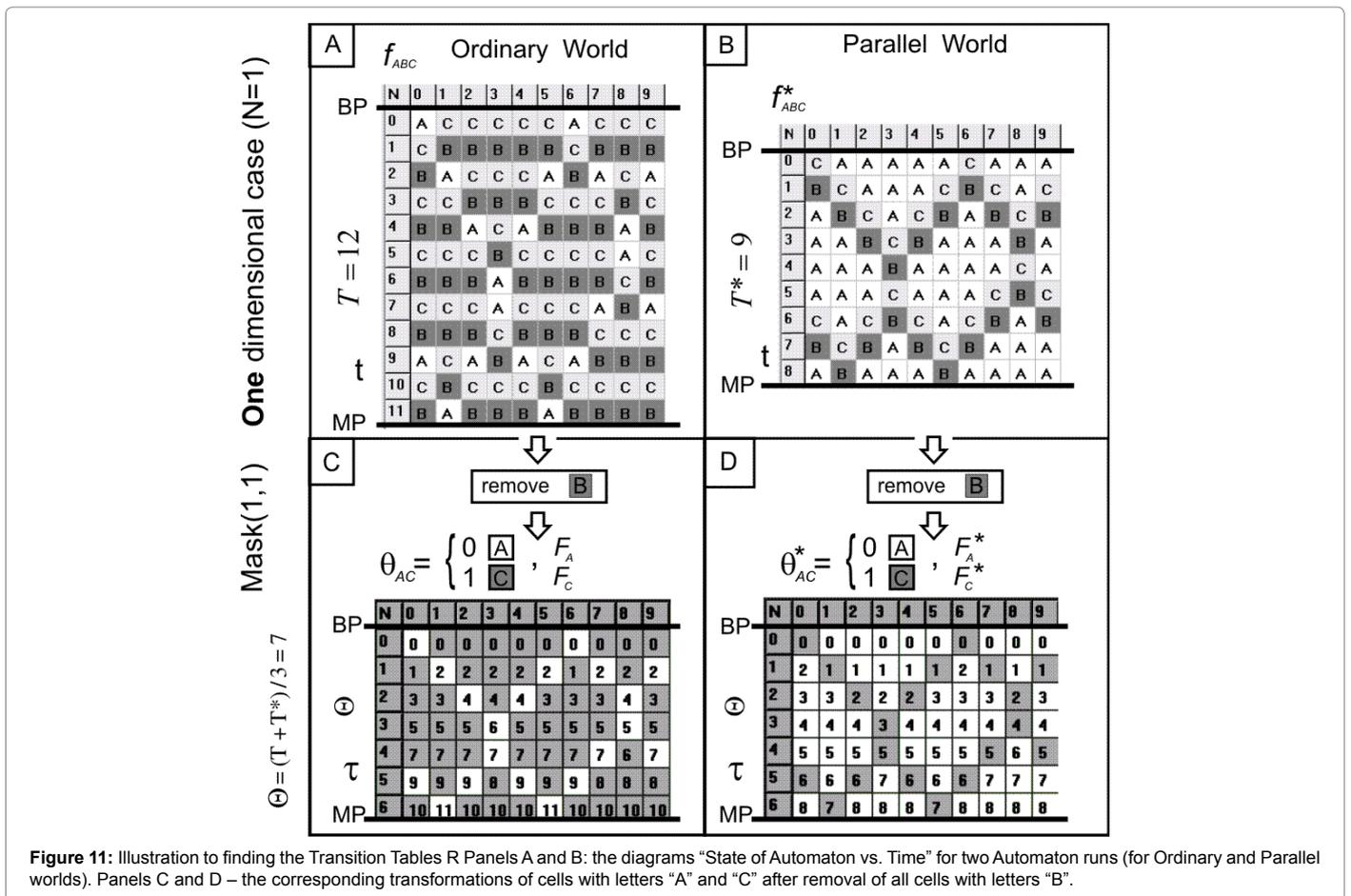
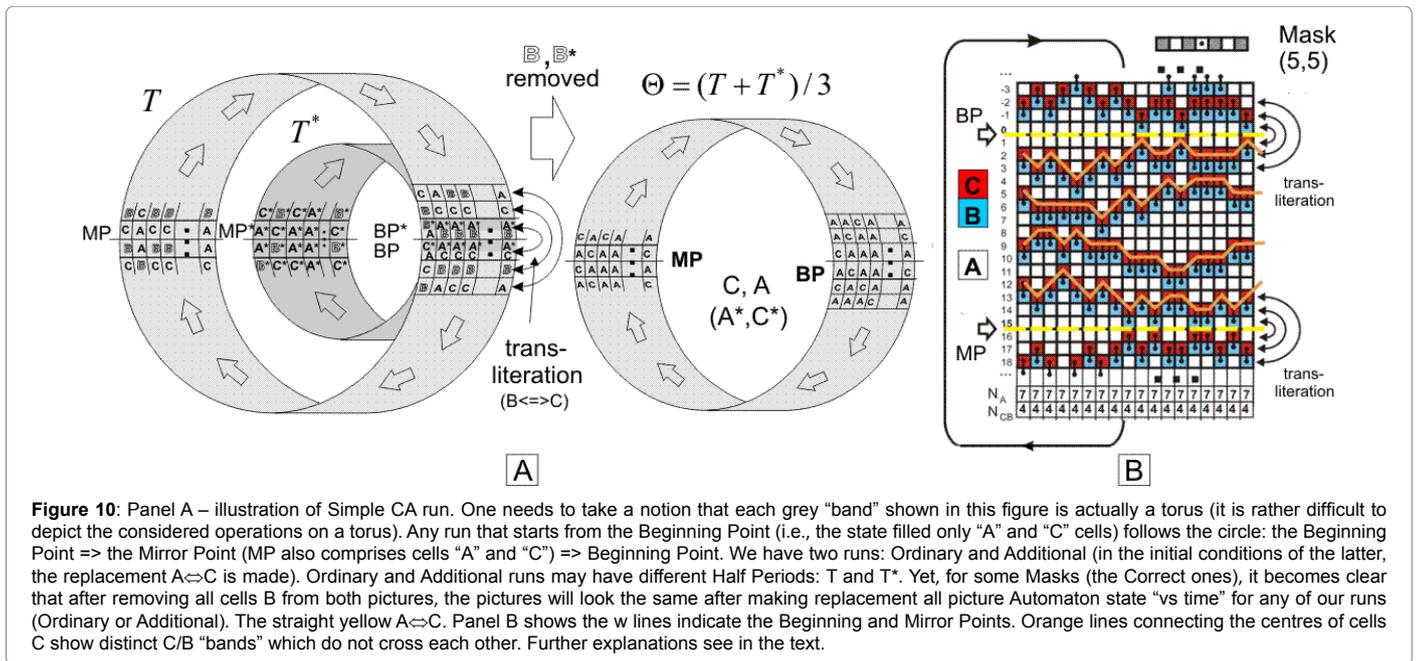
Let us return to the question how one obtains the Transition Table. To answer this, let us briefly describe the results presented [2]. Let us introduce the concept of Simple reversible Automaton on a Mask. It works on the same lattice (Z^N), yet closed on all sides (because we require a limited number of states), and has 3 states A, B, C. Here we describe an ordinary second-order Automaton found by E. Fredkin that

was studied by many authors, instance [3-5]. Each consecutive state depends on whether the Mask superimposed with a given point contains state C or not. In Case 1 (when the Mask contains state C) the following transitions are made $A \Rightarrow C, B \Rightarrow A, C \Rightarrow B$. In opposite Case 2, (no state C is contained in the Mask) the following transitions are made: $A \Rightarrow A, B \Rightarrow C, C \Rightarrow B$. It is well known that Fredkin's Automaton is reversible [6]), thus to change the direction of time, it is enough to replace all B with C and all C with B. Let us fill the cells only with the states taken from A and C ($t=0$) and consider the of the Automaton runs from the Beginning (Start) Point to the Mirror Point, and then back to the Beginning Point (Figure 10). An additional commentary to Figure 10. A normality of the Mask is necessary for two reasons. The first reason is that the Beginning point (BP) comprising cells A and C should have the Mirror point (MP) which is also composed of cells A and C (Figure 10A). The second reason is that the run of Automaton on the plot "Automaton state vs. Time" follows the distinct and non-crossing bands (Figure 10B). We will remind, in brief, how it has been proven.

Imagine a chess figure which makes only a one-unit move forward or backward along any of axes of the N-dimensional integer lattice (at each step, the figure can move along some arbitrary axis); also note that it is permitted for this figure to make a move on the field where the figure was in the past. It is obvious that the figure can eventually visit all the points of any closed integer lattice and return to its starting position. So, using the course of this figure we transform any N dimensions into one as shown in Figure 10B). After that, to prove that our bands do not cross each other bands, we use the property of a normality, [2].

Further, there is the main observation that, for any Correct Mask, CA of arbitrary size, and any initial conditions selected in the diagram "Automaton state vs. Time" where all B cells are removed, pictures for the "Ordinary run" and "Additional run" (wherein the replacement of the initial conditions $A \Leftrightarrow C$ is made), literally coincide after the replacement all of cells $A \Leftrightarrow C$. Consider the issue again and in detail. Let us name the diagram "Automaton state vs. time" for the normal





run of Automaton as the Ordinary World; and a similar diagram for the “additional run”, as Parallel World. First, let us introduce the terms describing the Ordinary World. The state of Simple CA in point r as a

function of time t is denoted as

$$f_{ABC}(t,r) \in \{A;B;C\} \quad (7)$$

Then, let us introduce function b which indicates whether a given cell is the B cell at a given time, or not

$$b(t, r) = \begin{cases} 1, & \text{if } f_{ABC}(t, r) = B \\ 0, & \text{if } f_{ABC}(t, r) \neq B \end{cases} \quad (8)$$

Transformation functions (with B removed) are defined as follows also explained in Figure 11.

$$\theta_{AC}(t - \sum_{0 \leq i \leq t} b(i, r), r) = \begin{cases} 0, & \text{if } f_{ABC}(t, r) = A \\ 1, & \text{if } f_{ABC}(t, r) = C \end{cases}$$

$$F_A(t - \sum_{0 \leq i \leq t} b(i, r), r) = \begin{cases} -1, & \text{if } f_{ABC}(t, r) = C \\ t, & \text{if } f_{ABC}(t, r) = A \end{cases}$$

$$F_C(t - \sum_{0 \leq i \leq t} b(i, r), r) = \begin{cases} -1, & \text{if } f_{ABC}(t, r) = A \\ t, & \text{if } f_{ABC}(t, r) = C \end{cases}$$

Let us take the primitive Mask in one dimension on the closed row (a ring) of ten points, two points of 10 being colored with two different ("opposite") colors. Yet, we repeat one more time, that all our methods can be applied without any changes to any Correct Mask in the N-dimensional space, for Automaton of any size, and for any initial conditions (Figure 12).

Let us introduce four new functions as follows:

$$F_{AC}(\tau, r) = \begin{cases} F_A(\tau, r), & \text{if } \theta_{AC}(\tau, r) = 0 \\ F_C(\tau, r), & \text{if } \theta_{AC}(\tau, r) = 1 \end{cases} \quad F_{CA}(\tau, r) = \begin{cases} F_A(\tau, r), & \text{if } \theta_{AC}(\tau, r) = 1 \\ F_C(\tau, r), & \text{if } \theta_{AC}(\tau, r) = 0 \end{cases}$$

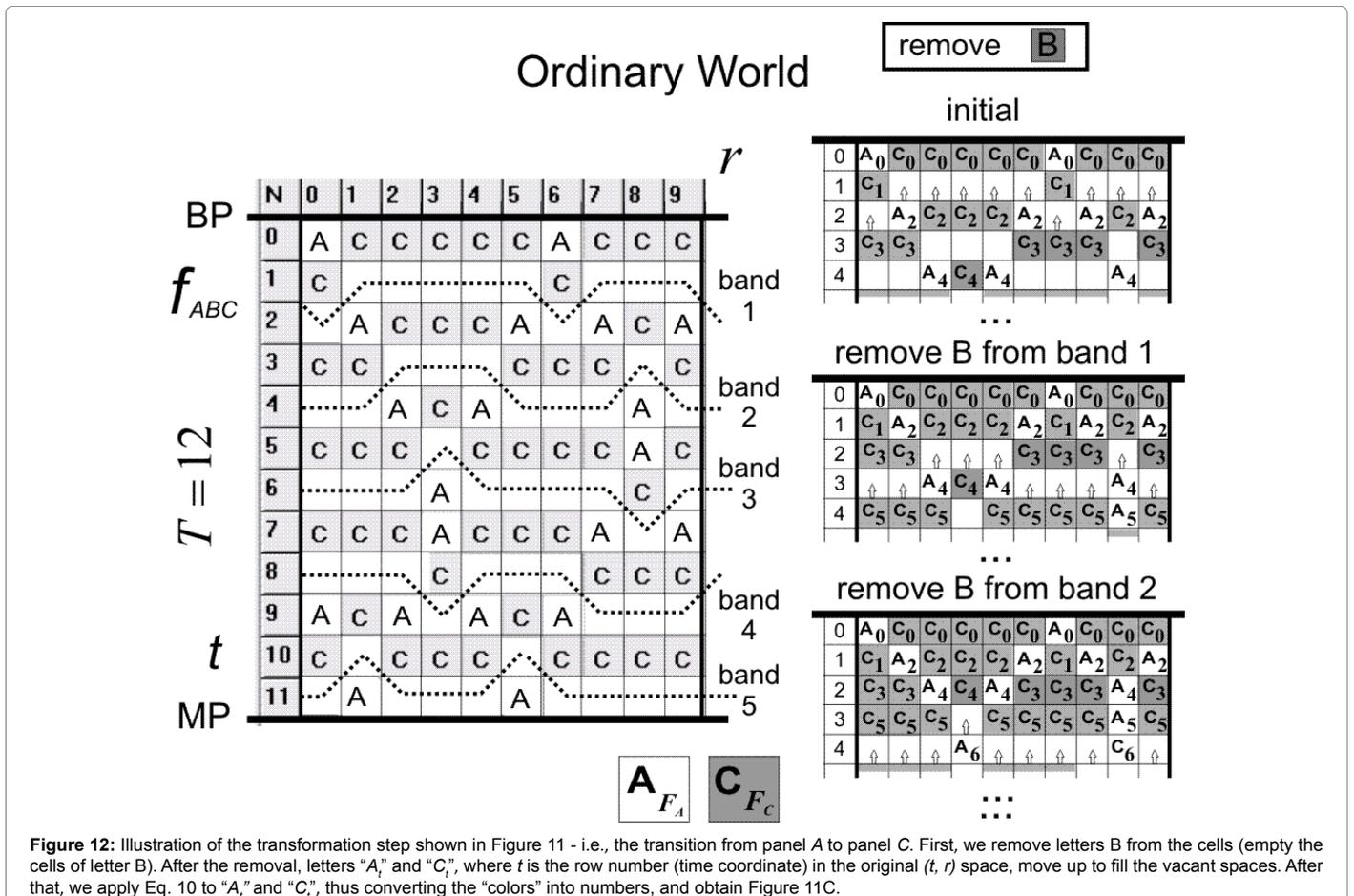
Note that the construction of $F_{AC}(\tau, r)$ was shown merely as an intermediate supplementary function to demonstrate the construction

of similar function $F_{AC}^{(3)}$ with the only difference in mod 3 (we emphasize that mod 3 is an important manipulation, it makes all other things happen). It should be understood that these operations create completely different tables $F_{AC}^{(3)}$ and $F_{CA}^{(3)}$. We introduce new notations: $0 \leftrightarrow -x, 1 \leftrightarrow -y, 2 \leftrightarrow -z, 3 \leftrightarrow +x, 4 \leftrightarrow +y, 5 \leftrightarrow +z$. We can see from the illustration of Figure 13 that these notations definitely give us the Transition Tables. The first Table is for one direction in time, and the second, for the opposite one. (If we have taken F_{CA} function instead of F_{AC} one, it would lead to replacement $R \leftrightarrow R^{-1}$).

Now it is clear how to obtain Transition Tables. First, we set random initial conditions from A and C states defined on "small" integer lattices that are closed in all directions and form the torus and turn on a Simple Automaton on the Mask. Then, for every three moves of a Simple Automaton on the time scale, we can find L (total cell number) multiplied by 2 states in each of our new tables $F_{AC}^{(3)}$. Finally, we stop after we reach the point of half-cycle and analyze the result. Then we change the initial conditions and repeat the trial. At each point in table $F_{AC}^{(3)}$, we can determine 12(!) new rows R_x taking into account the inverse motion. If the Mask is Correct, then R_x will be entirely filled in a short time. After that, it is necessary to perform standard check for Correctness (see Chapter 3).

Definition of a Perfect Mask. Strings

So, let us repeat again. Let us take a look on two runs (we call it as Standard and Additional) of i) a Simple Automaton on the Mask (f_{ABC} ; as a function of time t) and ii) Table Automaton on the same Mask $F_{AC}^{(3)}$



; as a function of time τ) induced by Simple Automaton. The cells we have obtained are the same and numbered with integers r . In the case of Simple Automaton, our run (the Ordinary World) starts with state filled with cells containing letters "A" and "C" at time $t=0$; Figure 10. In the additional run (Parallel World), the replacements $A \leftrightarrow C$ take place. In the case of Table Automaton, our run (the Ordinary World) starts with state filled $-x$ (cells with letters "A" in Simple Automaton; $t=0$) and $+x$ (cells with letters "C" in Simple Automaton; $t=0$). In the additional run (Parallel World) of the Table Automaton the replacements $-x \leftrightarrow x$ take place.

Let us introduce function $\Delta(\tau, r)$.

$$\Delta(\tau, r) = \begin{cases} -1 & \text{if } F_{AC}^{(3)}(\tau, r) = -x \vee -y \vee -z \\ +1 & \text{if } F_{AC}^{(3)}(\tau, r) = +x \vee +y \vee +z \end{cases}$$

And write out without any proof several formulas connecting our two Automaton. In any point r , parameters t and τ are connected by the relation:

$$t = \frac{3}{\tau} + \frac{1}{2} \sum_{0 \leq u \leq r} \Delta(u, r)$$

Eventually, our final observation is that if all this is correct for the Ordinary World, then for Parallel Word (the functions are designated by asterisk) and for all τ и r the following is true explained in Figure 14.

$$\sum_{0 \leq u \leq \tau} \Delta(u, r) = - \sum_{0 \leq u \leq \tau} \Delta^*(u, r) = 2d_{1/2}$$

$$\begin{aligned} \text{if } f_{ABC} \left(\frac{3\tau}{2} + d_{1/2}, r \right) = A &\Rightarrow f_{ABC}^* \left(\frac{3\tau}{2} - d_{1/2}, r \right) = C \\ \text{if } f_{ABC} \left(\frac{3\tau}{2} + d_{1/2}, r \right) = C &\Rightarrow f_{ABC}^* \left(\frac{3\tau}{2} - d_{1/2}, r \right) = A \end{aligned}$$

If for some Correct Mask and for every run the Eqs. 13-17 are satisfied, then we can call our Mask (as well as two of its CA) Perfect. Further on, we will portray this situation as "the two CA make their runs (from BP to the MP) by a use of a contiguous String". Note that theoretical physics has its own "strings", but even if the author has not known about it, he would have called this object with this word. We will show our String on some simple example (Figure 15). Let us show in the same drawing each of two one-dimensional runs of Simple Automaton for Masks (3,1). The length of row (closed into a ring) is 15 cells. Initial conditions are as follows: one point of "opposite" color (in Ordinary World: all cells C, one cell is A; in Parallel World: all cells A, one cell is C). Moreover, we use line of different colors to depict C/B «bands» for both runs: the green line indicate the C/B band in the Ordinary World while the blue line shows the C/B band in the Parallel World. Horizontal strips shown in panel A of Figure 15: x (light-blue) designate $t \bmod 3=0$, band y (white) $t \bmod 3=1$, and z (pink) $t \bmod 3=2$, where t is time (see above). Note that if we had taken normal but Incorrect Mask (5,1) instead of Mask (3,1), we would have obtained no Strings at all. Our beautiful picture from Figure 15 would have turned into a mess at once. Coming back to that chess figure from Section 4 which carries out converting the N dimensions of integer lattice into a single one, the author suggests (he didn't check it, yet) that whatever the method of

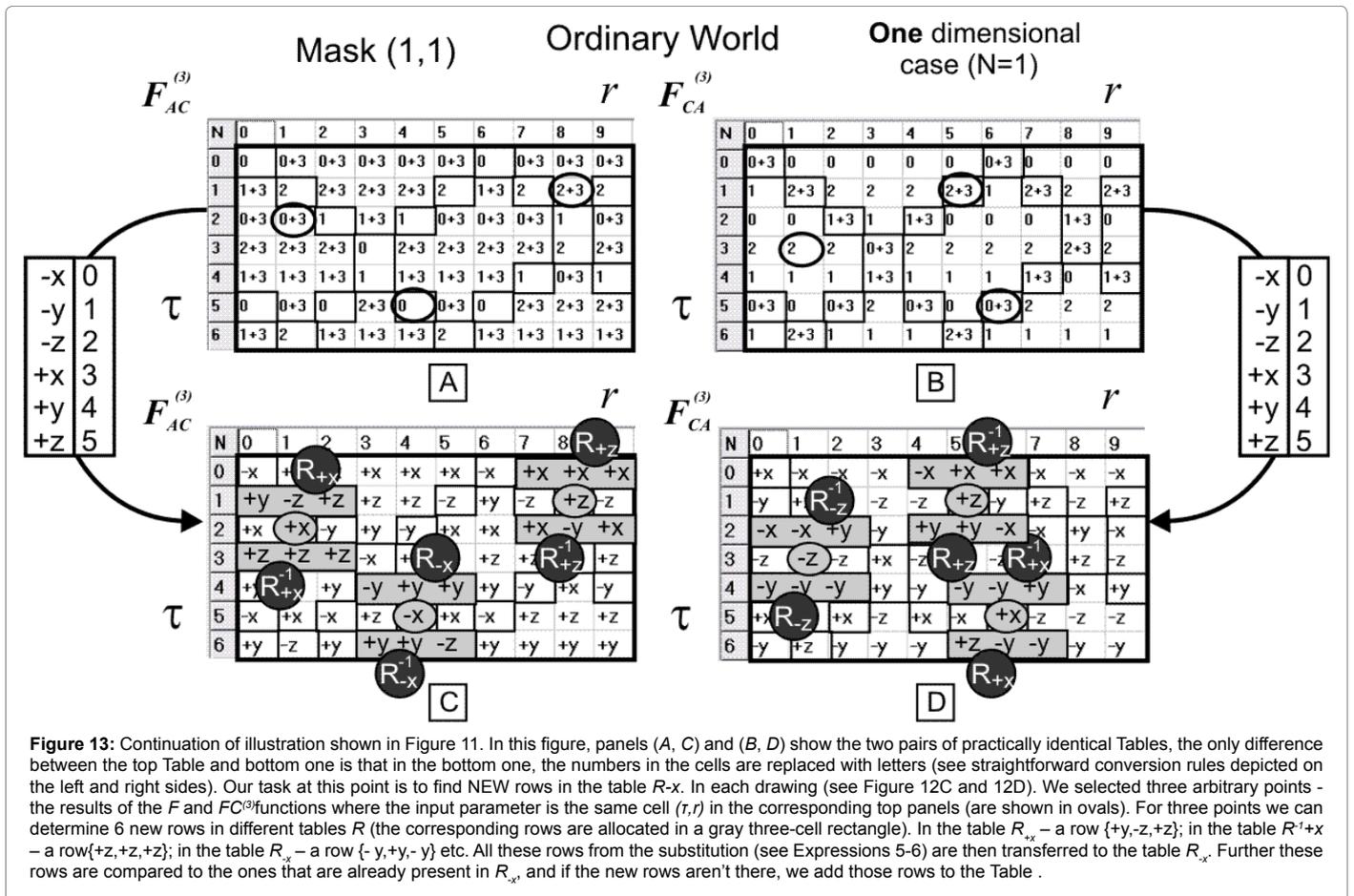
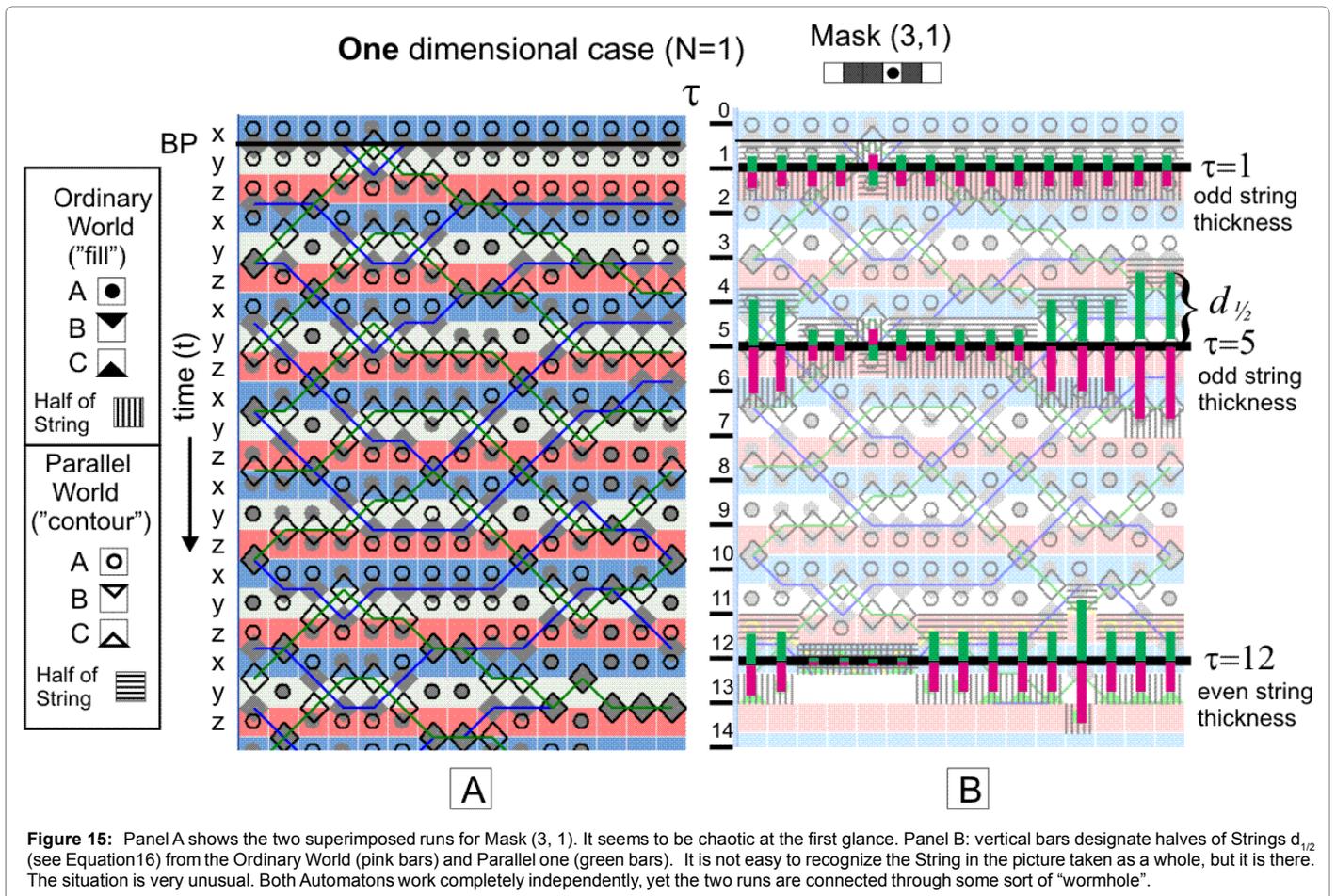
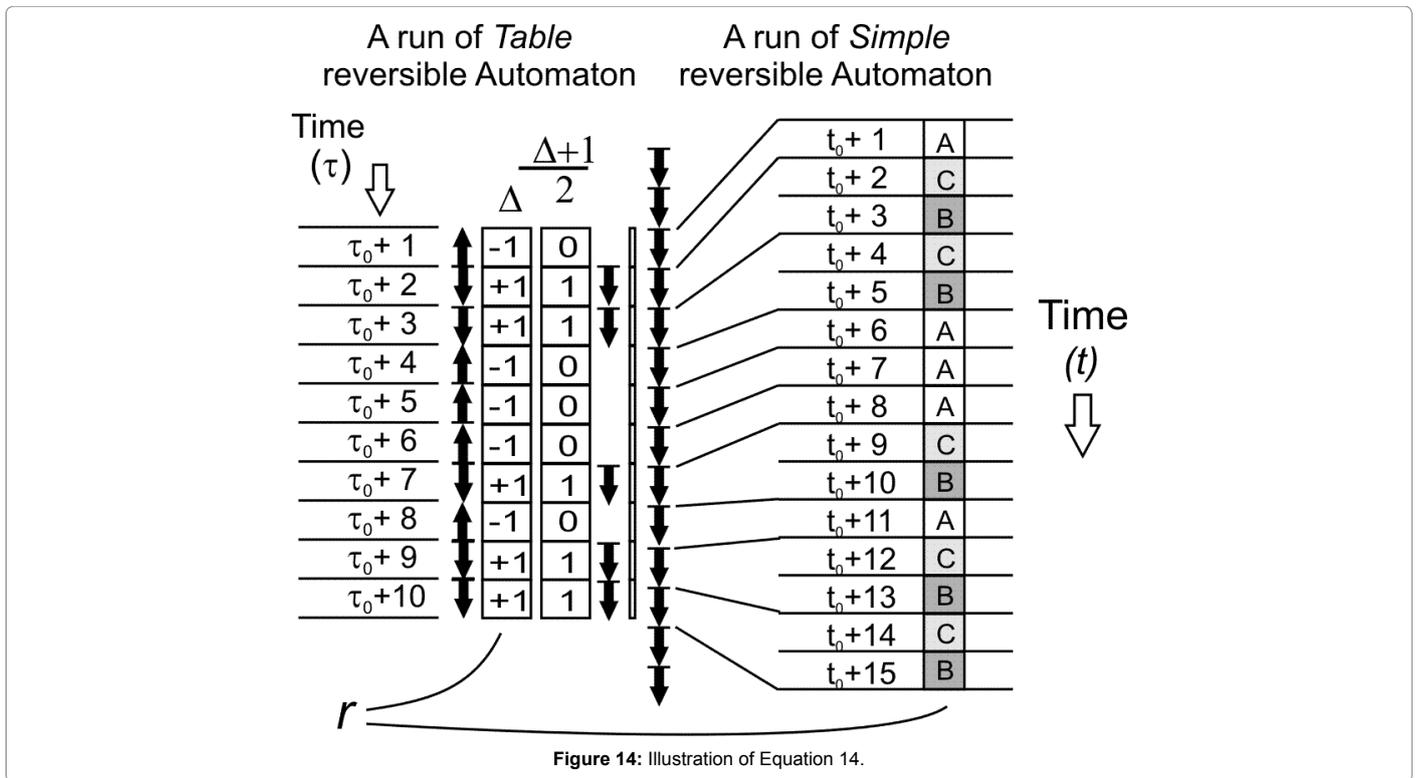


Figure 13: Continuation of illustration shown in Figure 11. In this figure, panels (A, C) and (B, D) show the two pairs of practically identical Tables, the only difference between the top Table and bottom one is that in the bottom one, the numbers in the cells are replaced with letters (see straightforward conversion rules depicted on the left and right sides). Our task at this point is to find NEW rows in the table R_x . In each drawing (see Figure 12C and 12D). We selected three arbitrary points - the results of the F and $FC^{(3)}$ functions where the input parameter is the same cell (τ, r) in the corresponding top panels (are shown in ovals). For three points we can determine 6 new rows in different tables R (the corresponding rows are allocated in a gray three-cell rectangle). In the table R_x - a row $\{+y, -z, +z\}$; in the table R^{+x} - a row $\{+z, +z, +z\}$; in the table R_x - a row $\{-y, +y, -y\}$ etc. All these rows from the substitution (see Expressions 5-6) are then transferred to the table R_x . Further these rows are compared to the ones that are already present in R_x , and if the new rows aren't there, we add those rows to the Table.



transformation is chosen, the String will keep the same appearance: i.e., the distance (taken at modulus) between the adjacent cells (in one measurement) will never be more than one.

So, there goes a reference run presenting a single step of Table CA. Over the course of this step, one and a half steps of Simple CA occur. The dashed shaded horizontal and vertical cells – representing a “half” of the Strings from the Ordinary World and the Parallel one – move down one cell (if there is no cell B beneath these cells) or two cells (if there is cell B beneath these cell, then one needs to jump over it). At every step and in every cell, the following event happens: either Ordinary World can go ahead of the Parallel World stays behind by the same amount) or lag (and then Parallel World is ahead). Parity of the String width ($d=2^*d_{1/2}$; the distance between the differently shaded cells) in each time point coincides with the parity of the number τ itself (Figure 15). At the Half Period point, Strings are stretched into a line with width $\lambda=T - T^*$ from [2]. In Section 4, we did not use the notion of Parallel World anywhere. In principle, the Correctness of Masks could exist without existence of Eqs. 13-17. But, it is obvious that these things are connected directly. If the Mask is Correct, that means there Eqs. 13-17 exist; if the Mask is Incorrect then Eqs. 13-17 do not exist. Naturally, the fact that the Mask is Perfect leads to a considerable reduction in the number of possible “ways” (rows of the states, parameter, Figure 6) and as a consequence, to a significant reduction (in a statistical sense) in the values of both half periods T and T* (see [7]). Immediately, there arises the Main problem (more correctly, a lot of problems). We denote it as the “Problem of the Number 3”. We do not know, yet, how broad is the problem. It is quite possible that it can be one of the main problems in Mathematics. Assume we know that the Mask under consideration is Correct. More accurately, we have found the Transition Tables with the aid of personal computer and proved by exhaustion that the given Mask is Correct. Now, we need to prove that the same very Mask is Perfect. Here arises a fundamental question of why the proof of Theorem 1 always works. In some sense, it can be the case that we only came up with a new and “better looking” wrap for the Correct Masks.

It is possible to suggest that analogously to Theorem 1 we can also prove the Eq. 16. Yet, what to do next? Well, it is necessary to investigate algebraically our Strings in N dimensions. Obviously, it is quite a challenging task.

How many “problems” do we have? In other words, how large is the number of all Correct Masks? We introduce, informally, certain number K which is the approximate number of Masks in all dimensions, for which the Jaguar - the most powerful supercomputer on this day (February 2016) - can prove the Correctness of all the Masks in one day or maximum a week. (This definition, obviously, is necessary to correct somehow. It is well known, for example, that in one dimension, all Masks ($-2^k+1, 2^k +1$) are Correct and C (=81) for those Masks is not very large. We, of course, do not consider such Masks! We do not know how large is K. It can be equal to 10,000 or may be equal to a factorial of 100. No doubts that the Jaguar is a powerful computer. On the other hand, the number of inspections grows very rapidly with the increase of n and C. However, the inspections use the algebraic operations/steps of similar type. So perhaps it is better to do the calculations using a specialized computer.

As of today, the general program has been written to prove the Correctness of any Mask in one and two dimensions. shows the proofs of the Correctness for all of Masks from Figure 6 with $n < 9$. The breadth of the proof at this point is limited by a performance of author’s personal computer. The Jaguar can increase “n” to several dozen (in one dimension) and number “N” to four or even five. In brief, Mathematical Strings definitely exist and in huge numbers. The author still does not

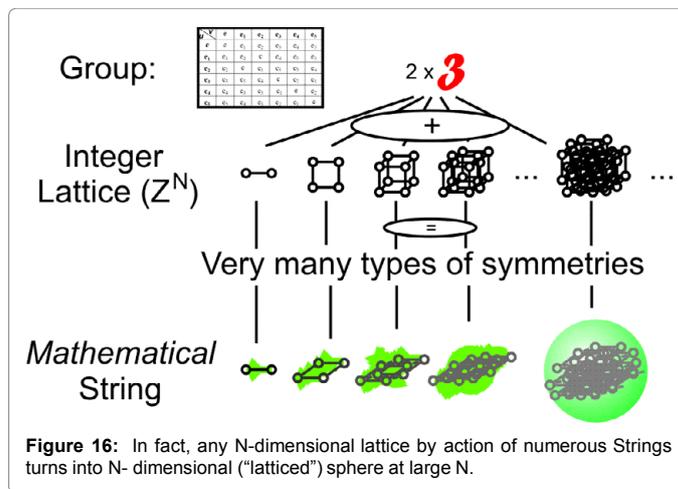


Figure 16: In fact, any N-dimensional lattice by action of numerous Strings turns into N- dimensional (“lattice”) sphere at large N.

know the answer to the question: whether or not all central-symmetric Masks are Correct Masks in all dimensions. The author employed the personal computer for weeks searching for abnormalities. The dimensions 1, 2, 3 have been checked. The criterion of Incorrectness was getting 2 in the expression $AC AC^*$ (Figure 11). Yet, not a single Incorrect Mask among central-symmetric ones has been found. To make the situation more intriguing, we note that there exist special Incorrect but “self-healing” Masks-Strings [8]. We cannot exclude a situation that starting with some n and N our proof stops working. Yet, the author tends to believe that a probability of such an outcome is extremely low. We suggest that all the material presented in the above to be named as the foundation of Theory of Mathematical Strings. Note: our theory should not be confused it with a well-known theory of strings in physics.

Conclusion

Let us briefly summarize our findings. The Number 3 (in other words, a group from Figure 5, or configuration 2x3) generates the unimaginable number of different symmetries in the N-dimensional integer lattice Z^N . Here we remind that the number of symmetries is greater than the number of central-symmetric Masks in the N-dimensional integer lattice Z^N . Computer calculations indicate (yet, indirectly) that in Z^N with large N almost any Mask is Correct, and each Correct Mask represents a new “symmetry” in this Z^N . The author developed the approach to finding the Correct Masks for low-dimensional Z^N . Utilizing the same approach (yet exploiting capabilities of a very powerful computer) it is feasible to find the Correct Masks for integer lattices of higher dimensions. Our main idea is illustrated by the following figure. It seems certain that these symmetries are directly related to the physical nature of space. We do this to suggest that professionals in the field of theory of numbers should take over the theoretical physicists concerning the question of constructing a consistent theory of the Universe. Yet, while physicists argue and debate about various peculiarities of string theory, one can contemplate how to use the theory of Mathematical Strings to build a consistent model describing fundamental properties of space (Figure 16).

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